

# 1.5

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## 1 1.5, p.21

1. (a)  $\frac{(\forall x)(\forall y)(x + y = 2)}{\text{free variables: none, so this is a sentence}}$   
 (b)  $\frac{(x + y < x) \vee (\forall z)(z < 0)}{\text{free variables: } x, y, \text{ so this is not a sentence}}$   
 (c)  $\frac{((\forall y)(y < x)) \vee ((\forall x)(x < y))}{\text{free variables: } x, y, \text{ so this is not a sentence}}$
6.  $\phi(x) := \underbrace{[(\forall y)(x = y)]}_{\alpha} \vee \underbrace{[(\forall x)(x < 0)]}_{\beta}, t \equiv S0$

The variable  $x$  is only free in  $\alpha$ , so  $\phi(t) := [(\forall y)(t = y)] \vee [(\forall x)(x < 0)]$

## 2 1.6, p.26

3.  $\mathcal{L}$  is  $\{b, \#^3, \natural^2\}$
7.  $\mathcal{L}_{NT}$  is  $\{0, S, +, \cdot, E, <\}$   
 $\mathfrak{N} = (\mathbb{N}, 0^{\mathfrak{N}}, S^{\mathfrak{N}}, +^{\mathfrak{N}}, \cdot^{\mathfrak{N}}, E^{\mathfrak{N}}, <^{\mathfrak{N}})$   
 $S^{\mathfrak{N}}(t) \equiv St$   
 $+^{\mathfrak{N}}(t, s) \equiv +ts$   
 $\cdot^{\mathfrak{N}}(t, s) \equiv \cdot ts$   
 $E^{\mathfrak{N}}(t, s) \equiv Ets$   
 $<^{\mathfrak{N}}(t, s) \equiv <ts$   
 $S0 + S0 \stackrel{?}{=} SS0$

### 3 1.7, p.32

1. The structure  $\mathfrak{N}$  makes the sentence  $1 + 1 = 2$  true.

*Proof.*



5. Let  $\mathfrak{A}$  be a structure for the language of set theory,  $\mathcal{L}_{ST}$ , which is  $\{\in\}$ .  
Let  $A = \{u, v, w, \{u\}, \{u, v\}, \{u, v, w\}\}$ .  
 $(\forall y \in y)(\exists x \in x)(x = y)$